

# The Paradoxical Success of Aspect- Oriented Programming

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# Introduction

- ❖ AOP sets out to increase modularity and structure of code by enabling the modularization of cross-cutting concerns
- ❖ AOP is a promising new technology; in many ways like OO once was
- ❖ AOP is being adopted by increasing numbers, both in industry and academia
- ❖ AOP works against independent development and understandability of code, two of the primary purposes of modularization
- ❖ Thus, AOP's success as a means of achieving modularization is paradoxical

# AOP

## - a Moving Target

- ❖ Each AOPL comes with its own (unambiguous) formal description of what AOP is
- ❖ No single definition that is
  - ~ common to all AOPLs and
  - ~ sufficiently distinguishes it from other, long established programming concepts
- ❖ There is though a common understanding what AOP is good for, namely modularizing cross-cutting concerns

# The Aspect Formula

- ❖ The (probably) best known definition AOP is

aspect-orientation = quantification +  
obliviousness

# Obliviousness

- ❖ Obliviousness means that a program has no knowledge of which aspects modify it or when
- ❖ Obliviousness as a defining characteristic of AOP has been questioned by the AOP community
- ❖ Some say that obliviousness is what distinguishes AOP from event-driven systems
- ❖ Obliviousness comes more as a side-effect of quantification

# Quantification

- ❖ Quantification means that an aspect can affect arbitrarily many different points in a program
- ❖ Quantification is widely accepted as a defining characteristic of AOP

# The Aspect Formula, cont'd

- ❖ The sentence “*In programs P, whenever condition C arises perform action A*” captures how an aspect (C, A) affects a given program P,
- ❖ but says nothing about P’s knowledge of the aspect (C, A), and thus nothing about obliviousness
- ❖ As the context provided to an action A is provided by the aspect (C, A) and not by the program P the program is oblivious to which program elements an aspect relies on, as opposed to a function call where arguments are explicitly passed to the function

# Interpretations of the Aspect Formula

- ❖ Translated in terms of AspectJ the parts of the formula read
  - ~ P is the execution of a program, which includes the execution of advice
  - ~ C is a set of pointcuts specifying the target elements of the aspect in the program and the context in which they occur (mostly variables, but also stack content)
  - ~ A is a piece of advice that depends on the context captured by C; and
  - ~ the quantification is implicit in the weaver



# Playing with the Formula

- ❖ Using different formulations of the condition  $C$  we can investigate AOP, or really the above definition

# Awareness Extreme

- ❖ Consider a condition  $C$  such as

In programs  $P$ , whenever an aspect is referenced, perform its associated action  $A$

- ❖ This expresses nothing more than the semantics of a standard procedure call
- ❖ This shows that quantification can be completely independent of obliviousness, as all places where condition  $C$  can arise are marked in the program text
- ❖ The programmer of  $P$  needs to know about which aspects are there, how they are named and how they work
- ❖ This is not AOP, but it shows that the “definition” of AOP is quite stretchable

# Obliviousness Extreme

- ❖ Consider a condition  $C$  such as

In programs  $P$ , whenever *Random* indicates it, perform action  $A$

- ❖ This means that all points in a program are implicitly marked, but execution of  $A$  remains uncertain
- ❖ The programmer of  $P$  may be aware of AOP, but has no knowledge of the existence or behavior of any aspect

# Taming Obliviousness

- ❖ The two previous examples are at the far extremes of the interpretation of the AOP formula
- ❖ There are, of course, less extreme interpretations of the formula

# Annotations

- ❖ Consider a condition such as

In programs  $P$ , whenever condition  $C$  arises where element  $B$  is referenced, perform action  $A$

- ❖  $B$  may be an abstract annotation
- ❖ Enables the programmer to deny aspects access where it is not wanted by not referencing  $B$ , but this means that the programmer must know of the aspects
- ❖ This is more or less equivalent to inserting a dynamically bound procedure call
- ❖ For massively crosscutting-concerns the annotations may very well turn out as annoying as the scattering of code that the aspect was to modularize

# Annotations, cont'd

- ❖ Using annotations reduces obliviousness to a level where the programmer of  $P$  knows that aspects may interact with the points marked  $B$  in  $P$ , but not which aspects or when
- ❖ However, annotations can act as interfaces between the program and the aspects, translating some of the best practices of OOP to AOP

# Annotations, cont'd

- ❖ Consider the following condition C

In programs P, whenever condition C arises, add annotation B

- ❖ Obviously the aspect could add the advice directly, but that would mean going back to the original formula

# Taming Quantification

- ❖ If and where aspects advice a program may very well seem random to a programmer
- ❖ Many propose tool support as a remedy to this, but tools can only mark the possible pointcut-“*shadows*” and not where and when advice are actually executed
- ❖ Keeping track of exactly where aspects advice an evolving program is not a trivial task
- ❖ One way of reducing this randomness is to use an explicit list of elements to be advised

In programs  $P$ , whenever execution reaches one of the points in  $\{p_1, \dots, p_n\}$ , perform action  $A$

- ❖ This is, of course, tedious and error prone for any interesting program



# Taming Quantification, cont'd

- ❖ Generally the quantification property of AOP suffers from the the problem that conditions are extremely sensitive to changes in the program (known as *the fragile pointcut problem*)
- ❖ Some researchers expect that this can be addressed by using better languages for expressing conditions, i.e. semantic pointcut languages.
- ❖ However, for an aspect to be useful in any interesting way it needs to reference the program context, at which point a semantic pointcut language cannot help,
- ❖ unless automatic program understanding is invented, which would revolutionize programming as a whole and render AOP, as well as every other technique known today, obsolete

# Modularity

- ❖ A module has a well defined interface which declares exactly what travels in and out of it
- ❖ This enables developers to work on different parts of a system (more or less) independently

# Modules and Interfaces

- ❖ Interfaces form the border between modules
- ❖ Interfaces represent the coupling between modules
  - ~ If the interface between two modules is empty, there is no coupling between them
- ❖ Interfaces should be made as explicit as possible to enable independent development

# AOP and Modularity

- ❖ AOP breaks the modularity of the program by modularizing cross-cutting concerns
- ❖ One could argue that this is for a good cause - and thus worth it
- ❖ What happens when cross-cutting concerns crosscut each other? And as soon as an aspect is woven it is part of the program and thus is a candidate for weaving of other aspects

# AOP and Modularity, cont'd

- ❖ Of course one could introduce annotations in the program to mark the places that should be advised by aspects, but this makes AOP no different from a subroutine call
- ❖ It also reintroduces the very scattering of a concern that AOP was to avoid

# Conclusion

- ❖ AOP sets out to modularize cross-cutting concerns, but its very nature breaks modularity
- ❖ It appears as this paradox cannot be resolved by tweaking the mechanics of AOP, because you end up with something which is very close to what we already have
- ❖ As a way of organizing code AOP does a good job by localizing a scattered concerns, but at the same time it breaks modularity of the program
- ❖ Thus, AOP's success as a means of achieving modularization is paradoxical